

is the only science which has made no progress since the revival of letters; logic is the only science which has produced no growth of symbols." In my view APL is in the best tradition of Boole, De Morgan, Jevons, and Venn.⁷³

One of the most striking features in Iverson's *A Programming Language* is his demonstration that "the generalized matrix product and the selection operations together provide an elegant formulation in several established areas of mathematics. A few examples will be chosen from two such areas, symbolic logic and matrix algebra."⁷⁴ Iverson proceeded to show how his notation leads to a natural extension of De Morgan's laws.⁷⁵

De Morgan's law:

$$A \wedge B \leftrightarrow \sim(\sim A) \vee \sim B$$

Iverson's extensions:

$$\wedge/U \leftrightarrow \sim\vee/\sim U$$

$$\neq/U \leftrightarrow \sim=/\sim U$$

In ordinary APL:

$$U \leftarrow ? 5 4 3 \rho 2$$

$$V \leftarrow ? 3 6 7 \rho 2$$

$$\wedge/, (U \neq. \wedge V) = \sim (\sim U) =. \vee (\sim V)$$

In J, the latest form of Iverson's notation, his 1962 example is executed as follows:

$$u = .?5 4 3 \$2$$

$$v = .?3 6 7 \$2$$

$$(u \sim: / .* . v) \sim: - . (-, u) = / . + . (-, v)$$

where:

$\sim:$ is NOT EQUAL; $*$ is AND; $\sim:$ is MATCH;
 $\sim.$ is NOT; and $+$ is OR.

In algebra a leading negative can be removed by changing the signs of all quantities in the expression that follows; in APL a leading NOT (\sim) can be removed by interchanging the pairs AND and OR, EQUALS and NOT-EQUALS, etc. In the following example both functions F and G remove redundant blanks from a string.

Ordinary APL:

$$\forall Z \leftarrow F \ S; U$$

$$[1] \quad Z \leftarrow (\sim U \wedge 1 \phi U \leftarrow S = ' \ ') / S$$

$$[2] \quad \nabla$$

$$\forall Z \leftarrow G \ S; U$$

$$[1] \quad Z \leftarrow (U \vee 1 \phi U \leftarrow S \neq ' \ ') / S$$

$$[2] \quad \nabla$$

Direct definition:

$$F: (\sim U \wedge 1 \phi U \leftarrow \omega = ' \ ') / \omega$$

$$G: (U \vee 1 \phi U \leftarrow \omega \neq ' \ ') / \omega$$

APL continues to grow in power, and Iverson's final example,⁷⁶ written but not executable as $+./$ in APL, can be executed in J as follows.

Given:

$$A = . 1 3 2 \theta . 2 1 \theta 1 . : 4 \theta \theta 2$$

$$B = . 4 1 . \theta 3 . \theta 2 . : 2 \theta$$

$$f = . \sim: \&\theta$$

$$h = . + / @ \# * \theta$$

Then:

$$(f A) + / . h B$$

$$4 6$$

$$6 4$$

$$6 1$$

Iverson's generalized matrix product found immediate application in his formal description of indexed addressing on the IBM 7090 computer,⁷⁷ which in one line made clear what takes half a page of text in the *Principles of Operation* manual for that machine. There are, of course, many similar examples in Reference 78.

Arrays and locative symbols

APL is often referred to as the *array processing language*, and its power does to a great extent come from its ability to work with arrays directly, a feature of increasing importance as vector processors and parallel computing become available. When we specify a place by giving its latitude and longitude, or define a point on a scatter diagram by giving its X and Y coordinates, we intend that two numbers should be taken together to identify one object. This is the first step in thinking in terms of what Sylvester called *multiple quantity*.

Stevinus was the first to show how forces combine in the manner we know as the parallelogram of